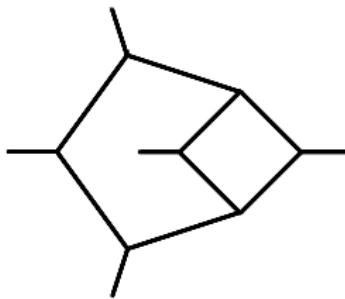


Integration-by-parts reductions via algebraic geometry

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University of Southampton



Amplitudes in the LHC era

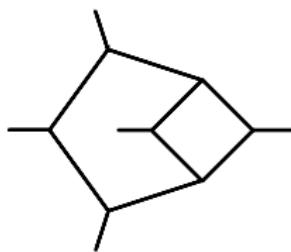
GGI Florence, 29th of October 2018

Based on PRD **93**(2016)041701, PRD **98**(2018)025023, JHEP **09**(2018)024
with J. Böhm, A. Georgoudis, H. Schönemann, M. Schulze, Y. Zhang

Overview

- 1 Motivation
- 2 IBP identities on unitarity cuts
- 3 Syzygy equations and their solution

- 4 Main example:



Integration-by-parts reductions

IBP identities arise from the vanishing integration of total derivatives,

[Chetyrkin, Tchakov, Nucl. Phys. B 192, 159 (1981)]

$$\int \prod_{i=1}^L \frac{d^D \ell_i}{\pi^{D/2}} \sum_{j=1}^L \frac{\partial}{\partial \ell_j^\mu} \frac{v_j^\mu P}{D_1^{a_1} \cdots D_k^{a_k}} = 0$$

where P and v_j^μ are polynomials in ℓ_i, p_j , and $a_i \in \mathbb{N}$.

Role in perturbative QFT calculations:

- **Reduction.** Reduce number of contributing loop integrals by factor of $\mathcal{O}(10^2) - \mathcal{O}(10^6)$ to basis.
- **Computing master integrals.** Enable setting up differential equations for basis integrals \mathcal{I}_j :

[Gehrmann and Remiddi, Nucl. Phys. B 580, 485 (2000)]

[Henn, PRL 110 (2013) 251601]

$$\frac{\partial}{\partial x_m} \mathcal{I}(\mathbf{x}, \epsilon) = A_m(\mathbf{x}, \epsilon) \mathcal{I}(\mathbf{x}, \epsilon)$$

where x_m denotes a kinematical invariant.

IBP reductions on unitarity cuts

Standard approach: enumerate all linear relations and apply
Gauss-Jordan elimination to *large* linear systems

[Laporta, Int.J.Mod.Phys. A 15 (2000) 5087-5159]

Idea here: use *unitarity cuts* to block-diagonalize system

We use the Baikov representation ($k = \frac{L(L+1)}{2} + LE$),

$$I(N; a) \equiv \int \prod_{j=1}^L \frac{d^D \ell_j}{i\pi^{D/2}} \frac{N}{D_1^{a_1} \cdots D_k^{a_k}} = \int \frac{dz_1 \cdots dz_k}{z_1^{a_1} \cdots z_k^{a_k}} \frac{\text{Gram}(z)^{\frac{D-L-E-1}{2}}}{(\hat{p}, \ell)} N$$

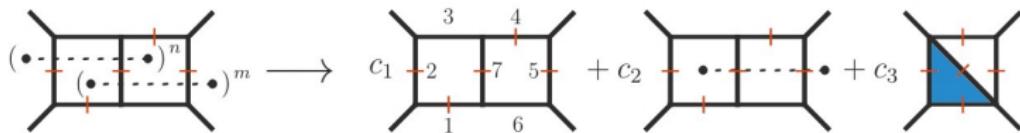
[Baikov, Phys.Lett. B 385 (1996) 404-410]

in which cuts are straightforward to apply,

$$\int \frac{dz_i}{z_i^{a_i}} \xrightarrow{\text{cut}} \oint_{\Gamma_\epsilon(0)} \frac{dz_i}{z_i^{a_i}} \quad i \in \mathcal{S}_{\text{cut}}$$

Example: Zurich-flag cut

Let us construct IBP identities on the Zurich-flag cut



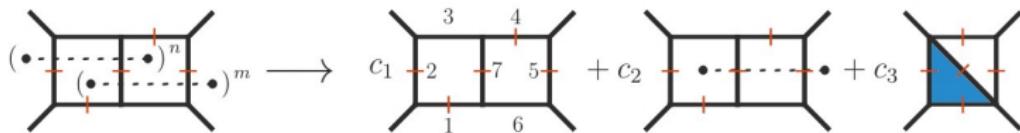
Define $S_{\text{cut}} = \{1, 2, 4, 5, 7\}$ and $G = \text{Gram}_{(\hat{\boldsymbol{p}}, \ell)}$.

On S_{cut} , the double-box integral takes the form

$$I_{\text{cut}}^{\text{DB}}[P] = \prod_{i \in S_{\text{cut}}} \oint_{\Gamma_\epsilon(0)} \frac{d\tilde{z}_i}{\tilde{z}_i} \int \prod_{j \notin S_{\text{cut}}} d\tilde{z}_j \frac{G(\tilde{\mathbf{z}})^{\frac{D-6}{2}}}{\tilde{z}_3 \tilde{z}_6} P(\tilde{\mathbf{z}})$$

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Relabeling $z_{\{1,2,3,4\}} = \tilde{z}_{\{3,6,8,9\}}$, this becomes

$$I_{\text{cut}}^{\text{DB}}[P] = \int \frac{dz_1 dz_2 dz_3 dz_4}{z_1 z_2} G(\mathbf{z})^{\frac{D-6}{2}} P(\mathbf{z})$$

Generic total derivative on cut

Need to find IBP identities which involve

$$I_{\text{cut}}^{\text{DB}}[P] = \int \frac{dz_1 dz_2 dz_3 dz_4}{z_1 z_2} G(z)^{\frac{D-6}{2}} P(z)$$

Total derivatives → **IBP identities.** Generic total derivative on cut:

$$\begin{aligned} 0 &= \int \left[\sum_{i=1}^4 \frac{\partial}{\partial z_i} \left(\frac{a_i(z) G(z)^{\frac{D-6}{2}}}{z_1 z_2} \right) \right] dz_1 \cdots dz_4 \\ &= \int \left[\sum_{i=1}^4 \left(\frac{\partial a_i}{\partial z_i} + \frac{D-6}{2G} a_i \frac{\partial G}{\partial z_i} \right) - \sum_{j=1,2} \frac{a_j}{z_j} \right] \frac{G(z)^{\frac{D-6}{2}}}{z_1 z_2} dz_1 \cdots dz_4 \end{aligned}$$

The **red term** corresponds to an integral in $(D - 2)$ dimensions, and the **purple term** in general produces doubled propagators.

IBP identities from syzygies

To avoid **dimension shifts** and **doubled propagators** in

$$0 = \int \left[\sum_{i=1}^4 \left(\frac{\partial a_i}{\partial z_i} + \frac{D-6}{2G} a_i \frac{\partial G}{\partial z_i} \right) - \sum_{j=1,2} \frac{a_j}{z_j} \right] \frac{G(z)^{\frac{D-6}{2}}}{z_1 z_2} dz_1 \cdots dz_4$$

we demand that each term is *polynomial*,

$$\sum_{i=1}^4 a_i \frac{\partial G}{\partial z_i} + bG = 0$$

$$a_j + b_j z_j = 0$$

with a_i, b_i, b polynomials in z . Such eqs. are known as **syzygy equations**.

[Gluza, Kajda, Kosower, PRD83(2011)045012], [Schabinger, JHEP01(2012)077], [Ita, PRD94(2016)116015]

Obtain IBPs by plugging (a_i, b) into the top equation.

Note: (qa_i, qb) is also a solution, for polynomial q .

Strategy to solve syzygy equations

Solve syzygy equations with c cuts

$$a_j + b_j z_j = 0, \quad j = 1, \dots, k-c \quad (1)$$

$$\sum_{j=1}^{m-c} \textcolor{red}{a}_j \frac{\partial G}{\partial z_k} + \textcolor{red}{b} G = 0 \quad (2)$$

as follows.

- 1) The generators of (1) are trivial:

$$\mathcal{M}_1 = \langle z_1 \mathbf{e}_1, \dots, z_k \mathbf{e}_k, \mathbf{e}_{k+1}, \dots, \mathbf{e}_m \rangle$$

- 2) Generators $\mathcal{M}_2 = \langle (\textcolor{red}{a}_1, \dots, \textcolor{red}{a}_m, \textcolor{red}{b}), \dots \rangle$ of (2) for the *off-shell case* $c = 0$ **can be explicitly found**:

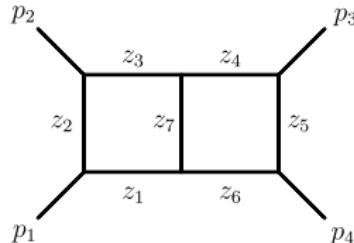
$$(\textcolor{red}{a}_\alpha, \textcolor{red}{b}) = \left(\sum_{k=1}^{E+L} (1+\delta_{ik}) x_{jk} \frac{\partial z_\alpha}{\partial x_{ik}}, \ 2\delta_{ij} \right)$$

where $x_{ij} = v_i \cdot v_j$ with $v_{i,j} \in \{p_1, \dots, p_E, \ell_1, \dots, \ell_L\}$.

[Böhm, Georgoudis, KJL, Schulze, Zhang, PRD 98(2018)025023]

- 3) Take module intersection $\mathcal{M}_1|_{\text{cut}} \cap \mathcal{M}_2|_{\text{cut}}$

Example 1: syzygies of planar double box



Set $P_{12} = p_1 + p_2$ and

$$\begin{aligned} z_1 &= \ell_1^2, & z_2 &= (\ell_1 - p_1)^2, & z_3 &= (\ell_1 - P_{12})^2 \\ z_4 &= (\ell_2 + P_{12})^2, & z_5 &= (\ell_2 - p_4)^2, & z_6 &= \ell_2^2 \\ z_7 &= (\ell_1 + \ell_2)^2, & z_8 &= (\ell_1 + p_4)^2, & z_9 &= (\ell_2 + p_1)^2 \end{aligned}$$

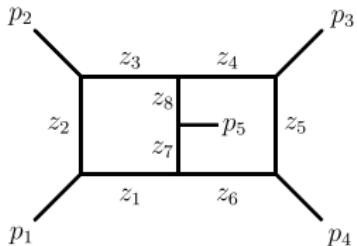
Only need to find explicit relation $z = Ax + B$. Here

$$A = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ -2 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ -2 & 0 & -2 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 2 & 0 & 2 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 2 & 0 & 2 & 0 & 2 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 2 & 1 \\ -2 & 0 & -2 & 0 & -2 & 0 & 1 & 0 & 0 & 0 \\ 0 & 2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

Set $t_{i,j} = (a_\alpha, b)$. The syzygy generators are *linear* in the z_k

$$\begin{aligned} t_{4,1} &= (z_1 - z_2, z_1 - z_2, -s + z_1 - z_2, 0, 0, 0, z_1 - z_2 - z_6 + z_9, t + z_1 - z_2, 0, 0) \\ t_{4,2} &= (s + z_2 - z_3, z_2 - z_3, z_2 - z_3, 0, 0, 0, z_2 - z_3 + z_4 - z_9, -t + z_2 - z_3, 0, 0) \\ t_{4,3} &= (-s + z_3 - z_8, t + z_3 - z_8, z_3 - z_8, 0, 0, 0, z_3 - z_4 + z_5 - z_8, z_3 - z_8, 0, 0) \\ t_{4,4} &= (2z_1, z_1 + z_2, -s + z_1 + z_3, 0, 0, 0, z_1 - z_6 + z_7, z_1 + z_8, 0, -2) \\ t_{4,5} &= (-z_1 - z_6 + z_7, -z_1 + z_7 - z_9, s - z_1 - z_4 + z_7, 0, 0, 0, -z_1 + z_6 + z_7, -z_1 - z_5 + z_7, 0, 0) \\ t_{5,1} &= (0, 0, 0, s - z_6 + z_9, -t - z_6 + z_9, z_9 - z_6, z_1 - z_2 - z_6 + z_9, 0, z_9 - z_6, 0) \\ t_{5,2} &= (0, 0, 0, z_4 - z_9, t + z_4 - z_9, -s + z_4 - z_9, z_2 - z_3 + z_4 - z_9, 0, z_4 - z_9, 0) \\ t_{5,3} &= (0, 0, 0, z_5 - z_4, z_5 - z_4, s - z_4 + z_5, z_3 - z_4 + z_5 - z_8, 0, -t - z_4 + z_5, 0) \\ t_{5,4} &= (0, 0, 0, s - z_3 - z_6 + z_7, -z_6 + z_7 - z_8, -z_1 - z_6 + z_7, z_1 - z_6 + z_7, 0, -z_2 - z_6 + z_7, 0) \\ t_{5,5} &= (0, 0, 0, -s + z_4 + z_6, z_5 + z_6, 2z_6, -z_1 + z_6 + z_7, 0, z_6 + z_9, -2) \end{aligned}$$

Example 2: syzygies of non-planar double pentagon



Set $P_{i,j} \equiv p_i + p_j$ and

$$\begin{aligned} z_1 &= \ell_1^2, & z_2 &= (\ell_1 - p_1)^2, & z_3 &= (\ell_1 - P_{1,2})^2, \\ z_4 &= (\ell_2 - P_{3,4})^2, & z_5 &= (\ell_2 - p_4)^2, & z_6 &= \ell_2^2, \\ z_7 &= (\ell_1 + \ell_2)^2, & z_8 &= (\ell_1 + \ell_2 + p_5)^2, & z_9 &= (\ell_1 + p_3)^2, \\ z_{10} &= (\ell_1 + p_4)^2, & z_{11} &= (\ell_2 + p_1)^2 \end{aligned}$$

Here $\mathbf{z} = \mathbf{Ax} + \mathbf{B}$ with

$$\mathbf{A} = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ -2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ -2 & 0 & 0 & -2 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & -2 & 0 & -2 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & -2 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 2 \\ -2 & -2 & -2 & -2 & -2 & -2 & -2 & -2 & -2 & 1 & 2 \\ 0 & 0 & 0 & 0 & 2 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 2 & 0 & 0 & 1 & 0 \\ 0 & 2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix},$$

$$\begin{aligned} t_{5,1} &= (z_1 - z_2, z_1 - z_2, -s_{1,2} + z_5 - z_2, 0, 0, 0, \\ &\quad z_1 - z_2 - z_6 + z_{11}, -s_{1,2} - s_{1,4} - s_{1,6} + z_1 - z_2 - z_6 + z_{11}, \\ &\quad s_{1,3} + z_1 - z_2, s_{1,4} + z_1 - z_2, 0, 0), \\ t_{5,2} &= (s_{1,2} + z_2 - z_3, z_2 - z_3, 0, 0, 0, \\ &\quad -s_{3,4} + z_1 + z_2 + z_4 + z_7 - z_8 - z_9 - z_{10} - z_{11}, \\ &\quad s_{1,3} + s_{1,4} + z_1 + z_2 + z_4 + z_7 - z_8 - z_9 - z_{10} - z_{11}, \\ &\quad s_{1,2} + s_{2,3} + z_2 - z_3, -s_{1,3} - s_{1,4} - s_{2,3} - z_3 + z_4 + z_5 + z_6, 0, 0), \\ t_{5,3} &= (z_0 - z_1, -s_{1,3} - z_1 + z_9, -s_{1,3} - s_{2,3} - z_1 + z_9, 0, 0, 0, \\ &\quad s_{3,4} - z_1 - z_4 + z_5 + z_9, -s_{1,3} - s_{2,3} - z_1 - z_4 + z_5 + z_9, \\ &\quad z_9 - z_1, s_{3,4} - z_1 + z_9, 0, 0), \\ t_{5,4} &= (z_{10} - z_1, -s_{1,4} - z_1 + z_{10}, -s_{1,4} - s_{2,4} - z_1 + z_{10}, \\ &\quad 0, 0, 0, -z_1 - z_5 + z_6 + z_{10}, \\ &\quad s_{1,2} + s_{1,3} + s_{2,3} - z_1 - z_5 + z_6 + z_{10}, \\ &\quad s_{3,4} - z_1 + z_{10}, z_1 - z_{10} - z_1, 0, 0), \\ t_{5,5} &= (2z_1, z_1 + z_2, -s_{1,2} + z_1 + z_8, 0, 0, 0, \\ &\quad z_1 - z_6 + z_7, -s_{1,2} + z_2) + z_8 - z_6 + z_7 - z_9 - z_{10}, \\ &\quad z_1 + z_9, z_1 + z_{10}, 0, -2), \\ t_{5,6} &= (-z_1 - z_6 + z_7, -z_1 + z_7 - z_3, \\ &\quad s_{1,2} + s_{3,4} - z_{21}, z_3 - z_4 + z_5 + z_9 + z_{10}, 0, 0, 0, \\ &\quad -z_1 + z_2 + z_7, s_{1,2} - z_{23}, z_3 + z_4 + z_5 + z_9 + z_{10}, \\ &\quad s_{3,4} - z_1 - z_4 + z_5 - z_6 + z_7, -z_1 - z_5 + z_7, 0, 0), \quad (6.11) \\ t_{6,1} &= (0, 0, 0, -s_{1,3} - s_{1,4} - z_1 + z_{11}, -s_{1,4} - z_6 + z_{11}, \\ &\quad z_{11} - z_6, z_1 - z_2 - z_6 + z_{11}, \\ &\quad -s_{1,2} - s_{1,3} - s_{1,4} + z_1 - z_2 - z_6 + z_{11}, 0, 0, z_1 - z_6, 0), \\ t_{6,2} &= (0, 0, 0, s_{1,3} + s_{1,4} + z_1 + z_3 + z_4 + z_7 - z_8 - z_9 - z_{10} - z_{11}, \\ &\quad s_{1,2} + s_{1,4} + s_{2,3} + z_1 + z_3 + z_4 + z_7 - z_8 - z_9 - z_{10} - z_{11}, \\ &\quad -s_{1,2} - s_{1,3} + z_1 + z_3 + z_4 + z_7 - z_8 - z_9 - z_{10} - z_{11}, \\ &\quad -s_{1,4} + z_1 + z_2 + z_4 + z_7 - z_8 - z_9 - z_{10} - z_{11}, \\ &\quad s_{1,3} + s_{1,4} + z_1 + z_2 + z_4 + z_7 - z_8 - z_9 - z_{10} - z_{11}, 0, 0, \\ &\quad -s_{1,4} + z_1 + z_3 + z_4 + z_7 - z_8 - z_9 - z_{10} - z_{11}, 0), \\ t_{6,3} &= (0, 0, 0, z_5 - z_4, z_5 - z_4, s_{3,4} - z_4 + z_5, \\ &\quad s_{3,4} - z_1 - z_4 + z_5 + z_9, -s_{1,3} - s_{2,3} - z_1 - z_4 + z_5 + z_9, 0, 0, \\ &\quad s_{1,3} - s_{3,4} - z_4 + z_5, 0), \\ t_{6,4} &= (0, 0, 0, -s_{3,4} - z_5 + z_6, z_6 - z_5, z_6 - z_5, \\ &\quad -z_1 - z_5 + z_6 + z_{10}, s_{1,2} + s_{1,3} + s_{2,3} - z_1 - z_5 + z_6 + z_{10}, 0, \\ &\quad s_{1,4} - z_5 + z_6, 0), \\ t_{6,5} &= (0, 0, 0, z_1 - z_6 + z_7 - z_9 - z_{10}, -z_6 + z_7 - z_{10}, \\ &\quad -z_1 - z_6 + z_7, z_1 - z_6 + z_7, -s_{1,2} + 2z_1 + z_3 - z_6 + z_7 - z_9 - z_{10}, \\ &\quad 0, 0, -z_2 - z_6 + z_7, 0), \\ t_{6,6} &= (0, 0, 0, -s_{3,4} + z_4 + z_6, z_6 + z_{10}, z_{10} - z_1 + z_6 + z_7, \\ &\quad z_{1,2} - z_1 - z_3 + z_6 + z_8 + z_9 + z_{10}, 0, 0, z_6 + z_{11}, -2), \end{aligned}$$

and the syzygy generators are again compact:

Computing module intersections

Given $\mathcal{M}_1 = \langle v_1, \dots, v_p \rangle$ and $\mathcal{M}_2 = \langle w_1, \dots, w_q \rangle$ with v_i, w_j m -tuples of polynomials. Let Q denote the $m \times (p+q)$ matrix

$$Q = \begin{pmatrix} \vdots & & \vdots & \vdots & & \vdots \\ v_1 & \cdots & v_p & w_1 & \cdots & w_q \\ \vdots & & \vdots & \vdots & & \vdots \end{pmatrix}$$

Then compute wrt. POT and variable order $[z_1, \dots, z_m] \succ [s_{ij}]$

$$\langle h_1, \dots, h_t \rangle \equiv \text{Gröbner basis of column space of } \left(\begin{array}{c|cc} Q & & \\ \hline 1 & & 0 \\ & \ddots & \\ 0 & & 1 \end{array} \right)$$

Selecting $h_i = (\overbrace{0, \dots, 0}^m, x_1, \dots, x_p, y_1, \dots, y_q)$, we have

$$0 = \sum_{j=1}^p x_j v_j + \sum_{k=1}^q y_k w_k$$

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$$Q = \begin{pmatrix} \vdots & & \vdots & \vdots & & \vdots \\ v_1 & \cdots & v_p & w_1 & \cdots & w_q \\ \vdots & & \vdots & \vdots & & \vdots \end{pmatrix}$$

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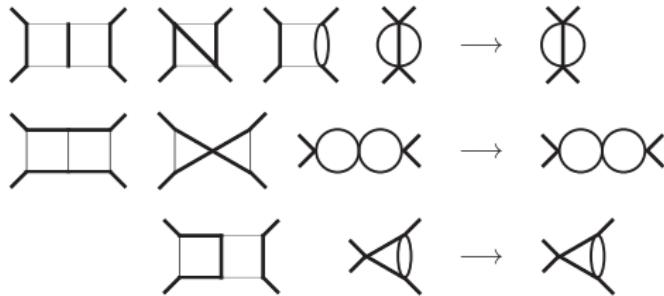
Selecting $h_i = (\overbrace{0, \dots, 0}^m, x_1, \dots, x_p, y_1, \dots, y_q)$, we have

$$0 = \sum_{j=1}^p x_j v_j + \sum_{k=1}^q y_k w_k \implies \sum_{j=1}^p x_j v_j = - \sum_{k=1}^q y_k w_k \in \mathcal{M}_1 \cap \mathcal{M}_2$$

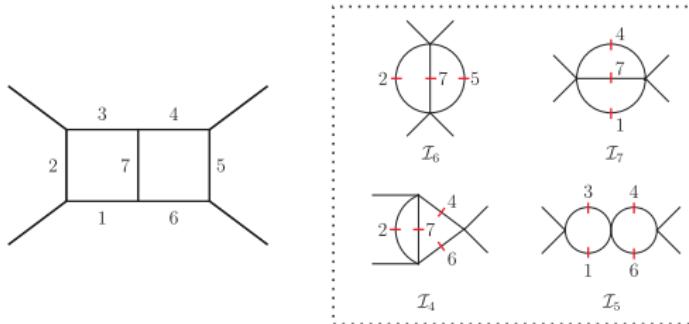
Hence $\sum_{j=1}^p x_j v_j$ generate $\mathcal{M}_1 \cap \mathcal{M}_2$, taking (x_1, \dots, x_p) from each h_i .

Spanning set of cuts for IBPs

To find the complete IBP reduction, we must consider the cuts associated with “uncollapsible” masters:



A bit more explicitly, the cuts we need to consider are



Main example: non-planar hexagon box

Task: IBP reduce non-planar hexagon box with numerator insertions of degree four in the z_i

[Chicherin, Henn, Mitev JHEP 05(2018)164]

[Badger, Brønnum-Hansen, Hartanto, Peraro, PRL 120(2018)092001]

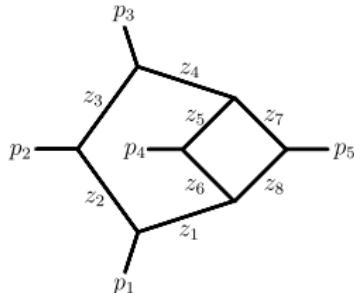
[Abreu, Cordero, Ita, Page, Zeng, PRD 97(2018)116014]

[Chawdhry, Lim, Mitov, 1805.09182]

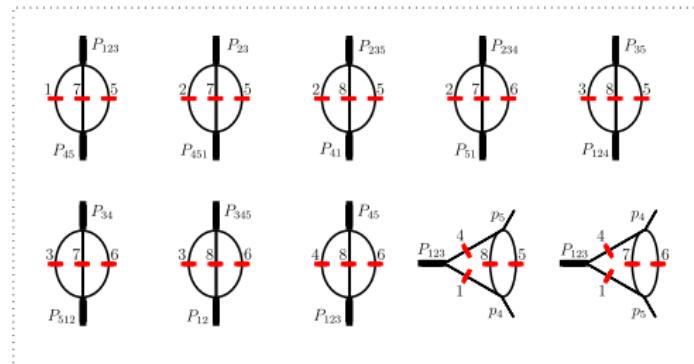
[S. Abreu, B. Page, M. Zeng, 1807.11522]

[D. Chicherin, T. Gehrmann, J. Henn,

N.A. Lo Presti, V. Mitev, P. Wasser, 1809.06240]

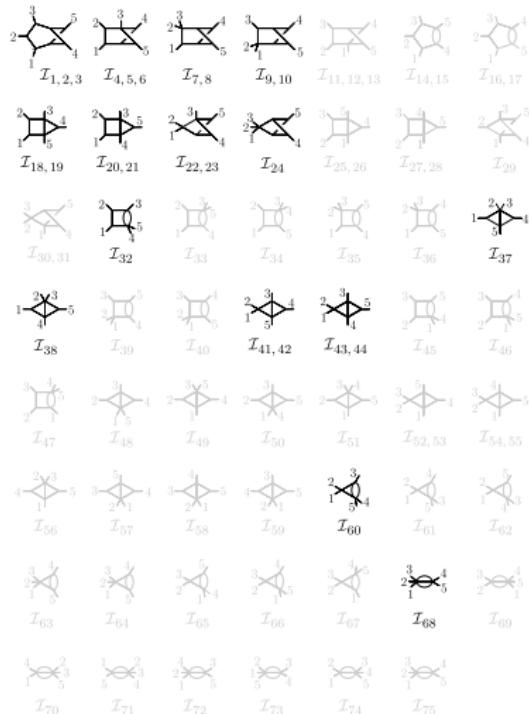


There are 10 cuts to consider:



Non-planar hexagon box: spanning set of cuts

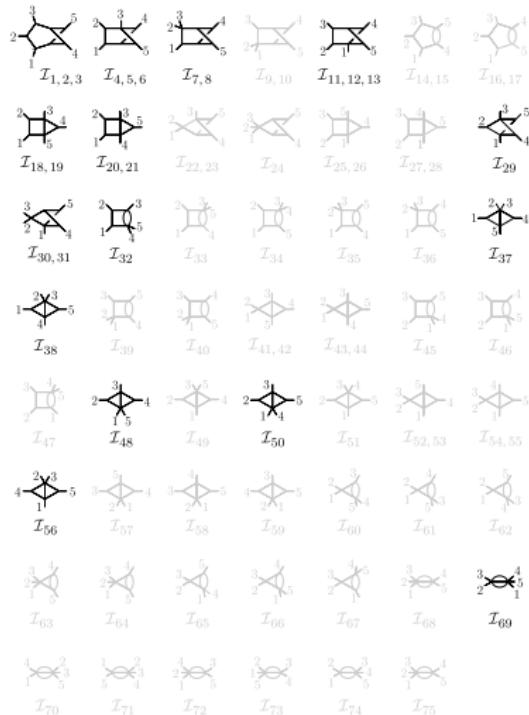
Construct and solve IBP identities on a spanning set of cuts.



Cut {1, 5, 7}

Non-planar hexagon box: spanning set of cuts

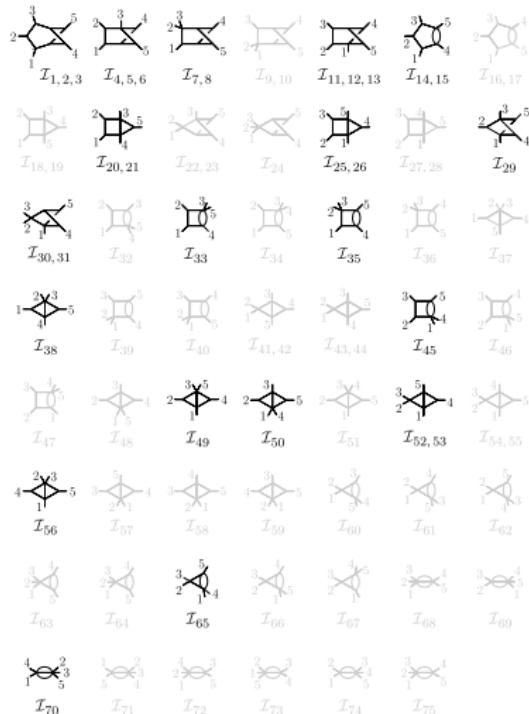
Construct and solve IBP identities on a spanning set of cuts.



Cut {2, 5, 7}

Non-planar hexagon box: spanning set of cuts

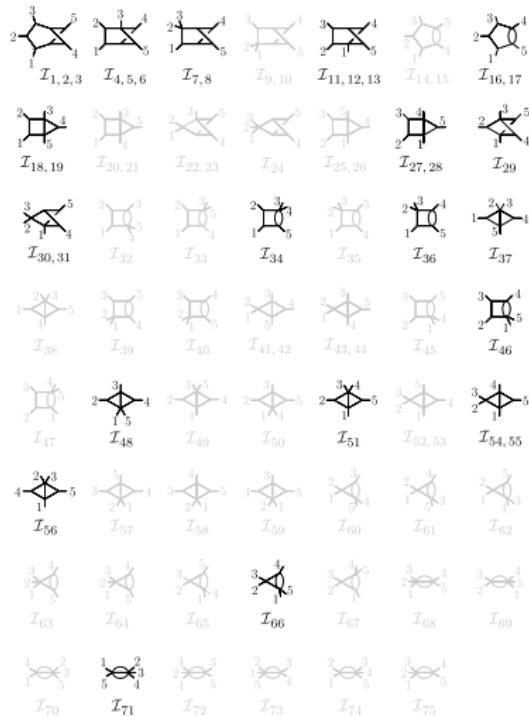
Construct and solve IBP identities on a spanning set of cuts.



Cut {2, 5, 8}

Non-planar hexagon box: spanning set of cuts

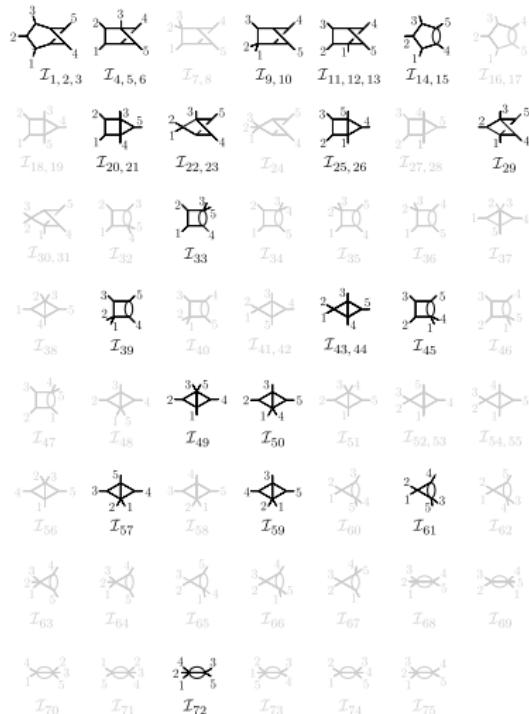
Construct and solve IBP identities on a spanning set of cuts.



Cut {2, 6, 7}

Non-planar hexagon box: spanning set of cuts

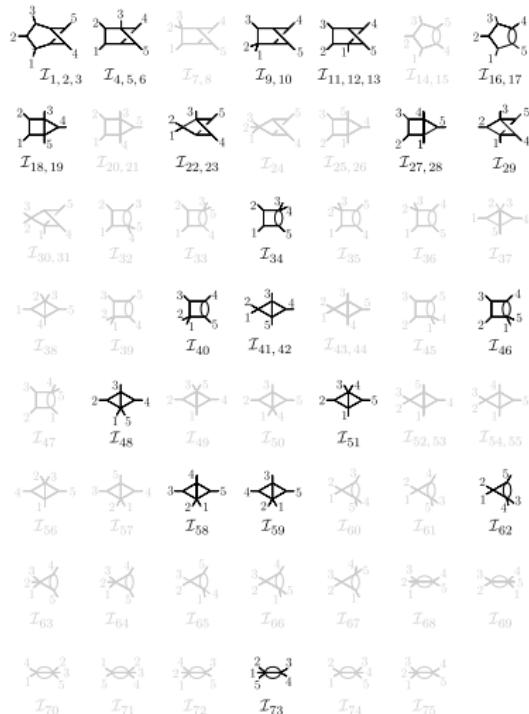
Construct and solve IBP identities on a spanning set of cuts.



Cut $\{3, 5, 8\}$

Non-planar hexagon box: spanning set of cuts

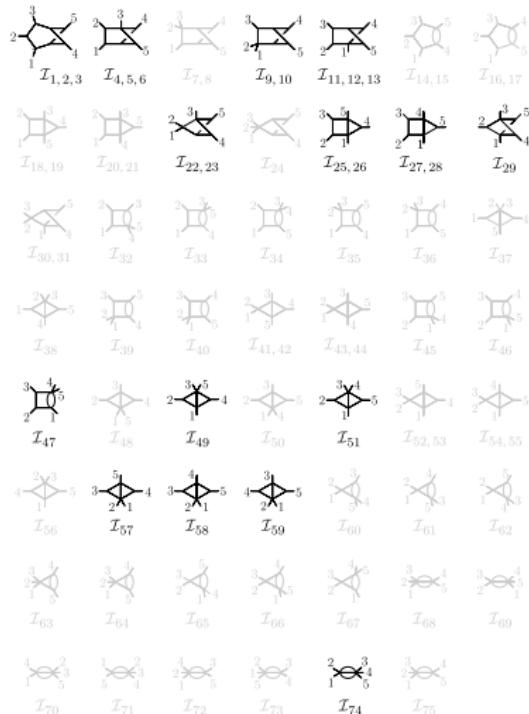
Construct and solve IBP identities on a spanning set of cuts.



Cut {3, 6, 7}

Non-planar hexagon box: spanning set of cuts

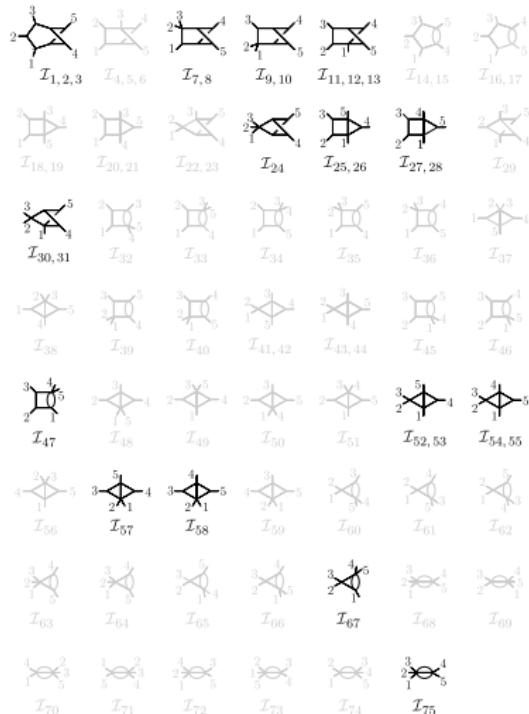
Construct and solve IBP identities on a spanning set of cuts.



Cut $\{3, 6, 8\}$

Non-planar hexagon box: spanning set of cuts

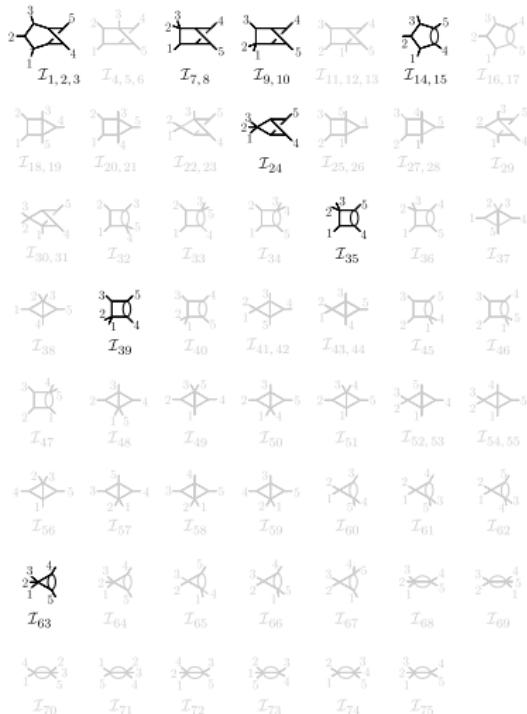
Construct and solve IBP identities on a spanning set of cuts.



Cut {4, 6, 8}

Non-planar hexagon box: spanning set of cuts

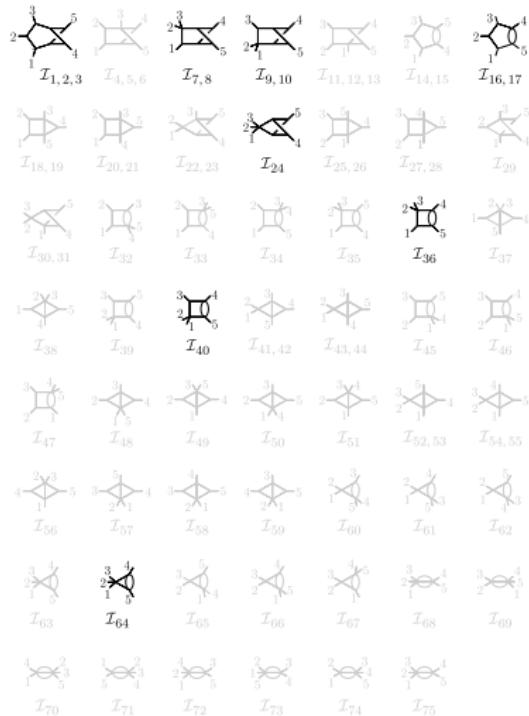
Construct and solve IBP identities on a spanning set of cuts.



Cut {1, 4, 5, 8}

Non-planar hexagon box: spanning set of cuts

Construct and solve IBP identities on a spanning set of cuts.



Cut {1, 4, 6, 7}

Syzygies for the non-planar hexagon box

Syzygies for ensuring D -dimensionality:

$$\begin{aligned}
M_1 = & \langle (z_1 - z_2, z_1 - z_2, -s_{12} + z_1 - z_2, -s_{12} - s_{13} + z_1 - z_2, s_{14} + z_1 - z_2 - z_8 + z_{10}, z_1 - z_2 - z_8 + z_{10}, 0, 0, -s_{12} - s_{13} - s_{14} + z_1 - z_2, 0, 0) \\
& (0, 0, 0, 0, s_{14} + z_1 - z_2 - z_8 + z_{10}, z_1 - z_2 - z_8 + z_{10}, s_{12} + s_{13} + s_{14} - z_8 + z_{10}, z_{10} - z_8, z_{10} - z_8, s_{12} - z_8 + z_{10}) \\
& (s_{12} + z_2 - z_3, z_2 - z_3, -s_{23} + z_2 - z_3, s_{12} + s_{24} + z_2 - z_3 - z_8 + z_{11}, s_{12} + z_2 - z_3 - z_8 + z_{11}, 0, 0, -s_{23} - s_{24} + z_2 - z_3, 0, 0) \\
& (0, 0, 0, 0, s_{12} + s_{24} + z_2 - z_3 - z_8 + z_{11}, s_{12} + z_2 - z_3 - z_8 + z_{11}, s_{12} + s_{23} + s_{24} - z_8 + z_{11}, z_{11} - z_8, 0, s_{12} - z_8 + z_{11}, z_{11} - z_8) \\
& (s_{13} + s_{23} + z_3 - z_4, s_{23} + z_3 - z_4, z_3 - z_4, z_3 - z_4, -2s_{12} - s_{13} - s_{14} - s_{23} - s_{24} + z_3 - z_4 + z_6 + z_7 + z_8 - z_9 - z_{10} - z_{11}, \\
& -s_{12} + z_3 + z_5 + z_6 + z_7 + z_8 - z_9 - z_{10} - z_{11}, 0, 0, s_{12} + s_{14} + s_{23} + s_{24} + z_3 - z_4, 0, 0) \\
& (0, 0, 0, 0, -2s_{12} - s_{13} - s_{14} - s_{23} - s_{24} + z_3 - z_5 + z_6 + z_7 + z_8 - z_9 - z_{10} - z_{11}, -s_{12} + z_3 - z_5 + z_6 + z_7 + z_8 - z_9 - z_{10} - z_{11}, \\
& -2s_{12} - s_{13} - s_{14} - s_{23} - s_{24} + z_4 - z_5 + z_6 + z_7 + z_8 - z_9 - z_{10} - z_{11}, -s_{12} - s_{13} - s_{23} + z_4 - z_5 + z_6 + z_7 + z_8 - z_9 - z_{10} - z_{11}, \\
& 0, -s_{12} - s_{23} + z_4 - z_5 + z_6 + z_7 + z_8 - z_9 - z_{10} - z_{11}, -s_{12} - s_{13} + z_4 - z_5 + z_6 + z_7 + z_8 - z_9 - z_{10} - z_{11}) \\
& (-s_{12} - s_{13} - s_{23} + z_4 - z_5 - z_6, -s_{12} - s_{13} - s_{14} - s_{23} + z_4 - z_5 - z_6, -s_{12} - s_{13} - s_{14} - s_{24} + z_4 - z_5 - z_6, z_5 - z_6, 0, 0, z_4 - z_5, 0, 0) \\
& (0, 0, 0, 0, z_5 - z_6, z_5 - z_6, -s_4 + z_5 - z_6 + z_9, s_{12} + s_{13} + s_{23} - z_4 + z_5 - z_6 + z_9) \\
& (s_{12} + s_{13} + s_{14} + s_{23} - z_4 + z_5 - z_6 + z_9, s_{12} + s_{13} + s_{23} + s_{24} - z_4 + z_5 - z_6 + z_9) \\
& (2z_1, z_1 + z_2, -s_{12} + z_1 + z_3, -s_{12} - s_{13} - s_{23} + z_1 + z_4, -s_{12} - s_{13} - s_{23} - s_{24} + z_1 + z_4 - z_9, z_4 - z_9, z_5 - z_6, z_5 - z_6, 0, 0, z_1 + z_9, 0, 0) \\
& (0, 0, 0, 0, -s_{12} - s_{13} - s_{23} + z_1 + z_4 - z_5 - z_6, -s_1 + z_1 + z_4 - z_5 - z_6, -z_6 - z_8 - z_9, -z_1 + z_6 - z_8, 0, -z_2 + z_6 - z_8, s_{12} - z_1 + z_2 - z_3 + z_6 - z_8) \\
& (-z_1 + z_6 - z_8, -z_1 + z_6 - z_10, -z_1 + z_6 + z_8 - z_{10} - z_{11}, s_{12} + s_{13} + s_{23} - z_1 - z_4 + z_5 - z_7 + z_9, \\
& z_{12} + s_{13} + s_{23} - z_1 - z_4 + z_5 + z_8 + z_9, -z_1 + z_6 + z_8, 0, 0, -z_1 + z_6 - z_7 + z_9, 0, 0) \\
& (0, 0, 0, 0, s_{12} + s_{13} + s_{23} - z_1 - z_4 + z_5 + z_8 + z_9, -z_1 + z_6 + z_8, z_7 + z_8, 2z_8, 0, z_8 + z_{10}, z_8 + z_{11}) \rangle \quad (5.9)
\end{aligned}$$

Syzygies for ensuring no doubled propagators:

$$\begin{aligned}
M_2 = & \langle (z_1, 0, 0, 0, 0, 0, 0, 0, 0, 0), (0, z_2, 0, 0, 0, 0, 0, 0, 0, 0) \\
& (0, 0, z_3, 0, 0, 0, 0, 0, 0, 0), (0, 0, 0, z_4, 0, 0, 0, 0, 0, 0) \\
& (0, 0, 0, 0, z_5, 0, 0, 0, 0, 0, 0), (0, 0, 0, 0, 0, z_6, 0, 0, 0, 0) \\
& (0, 0, 0, 0, 0, z_7, 0, 0, 0, 0), (0, 0, 0, 0, 0, 0, z_8, 0, 0, 0) \\
& (0, 0, 0, 0, 0, 0, 0, 1, 0, 0), (0, 0, 0, 0, 0, 0, 0, 0, 1, 0) \\
& (0, 0, 0, 0, 0, 0, 0, 0, 0, 1) \rangle
\end{aligned}$$

Compute intersection of $M_1|_{\text{cut}} \cap M_2|_{\text{cut}}$ on each of the 10 cuts.

Complexity of IBP systems

- Resources to compute $M_1|_{\text{cut}} \cap M_2|_{\text{cut}}$: **25-800 s** and **1-14 GB RAM**
(on 24 cores, 3.40 GHz)
- Size of generating systems after trimming: **1.5-10 MB**
Plug **resulting generators** into ansatz for total derivative:
$$0 = \int \left[\sum_{i=1}^{m-c} \left(\frac{\partial \mathbf{a}_{r_i}}{\partial z_{r_i}} + \frac{D-L-E-1}{2G(z)} \mathbf{a}_{r_i} \frac{\partial G}{\partial z_{r_i}} \right) - \sum_{i=1}^{k-c} \frac{\mathbf{a}_{r_i}}{z_{r_i}} \right] \frac{G(z)^{\frac{D-L-E-1}{2}}}{z_{r_1} \cdots z_{r_{k-c}}} dz_{r_1} \cdots dz_{r_{m-c}}$$
- Resulting linear systems to solve:
700-1200 equations, size **1 MB**, density **1.5%**

Gauss-Jordan elimination of IBP systems

To find the IBP reductions, Gauss-Jordan eliminate IBP systems.

Some remarks:

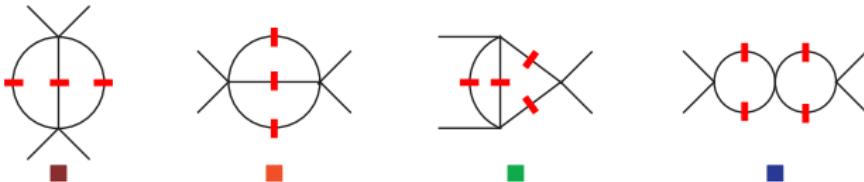
- To preserve sparsity, use a *total pivoting* strategy
(i.e., allow column swaps)
- For cut $\{1, 4, 6, 7\}$, the RREF can be performed fully analytically, requiring 31 minutes on one core and 1.5 GB RAM.
⋮
- For $\{3, 6, 7\}$, assigned numerical values to two s_{ij} .
Ran 440 points on cluster (2.5 h and 1.8 GB RAM per job).
Used interpolation code to get analytical results (23 min and 15 GB RAM on one core).

[von Manteuffel and Schabinger, PLB 744(2015)101]

[Peraro, JHEP12(2016)030]

Merging on-shell IBP reductions

By solving the IBP identities on the following cuts



we reconstruct the *complete IBP reductions* by merging the partial results.

An example of an IBP relation produced by our method ($\chi \equiv t/s$):

$$\begin{aligned} (\bullet \cdots \bullet)^2 &= \frac{(D-4)s^2\chi}{8(D-3)} \text{ (top-bottom cut)} - \frac{(3D-2\chi-12)s}{4(D-3)} \text{ (diagonal cut)} + \frac{(4-D)(9\chi+7)}{4(D-3)} \text{ (diagonal cut)} \\ &\quad + 2 \text{ (left-right cut)} + \frac{(10-3D)(2\chi-13)}{8(D-4)s} \text{ (top-bottom cut)} + \frac{2D(\chi+1)-8\chi-7}{2(D-4)s} \text{ (top-bottom cut)} \\ &\quad + \frac{9(3D-10)(3D-8)}{4(D-4)^2s^2\chi} \text{ (top-bottom cut)} + \frac{(3D-10)(3D-8)(2\chi+1)}{2(D-4)^2(D-3)s^2} \text{ (top-bottom cut)} \end{aligned}$$

Results for IBP reductions

- Fully analytic IBP reductions of the 32 hexagon boxes

$$\{ I(1, 1, 1, 1, 1, 1, 1, 1, 0, 0, -4), \quad I(1, 1, 1, 1, 1, 1, 1, 1, 0, -1, -3), \quad I(1, 1, 1, 1, 1, 1, 1, 1, 0, -2, -2) \\ I(1, 1, 1, 1, 1, 1, 1, 1, 0, -3, -1), \quad I(1, 1, 1, 1, 1, 1, 1, 1, 0, -4, 0), \quad I(1, 1, 1, 1, 1, 1, 1, 1, 1, -1, 0, -3) \\ I(1, 1, 1, 1, 1, 1, 1, 1, -1, -1, -2), \quad I(1, 1, 1, 1, 1, 1, 1, 1, -1, -2, -1), \quad I(1, 1, 1, 1, 1, 1, 1, 1, 1, -1, -3, 0) \\ I(1, 1, 1, 1, 1, 1, 1, 1, -2, 0, -2), \quad I(1, 1, 1, 1, 1, 1, 1, 1, -2, -1, -1), \quad I(1, 1, 1, 1, 1, 1, 1, 1, -2, -2, 0) \\ I(1, 1, 1, 1, 1, 1, 1, 1, -3, 0, -1), \quad I(1, 1, 1, 1, 1, 1, 1, 1, -3, -1, 0), \quad I(1, 1, 1, 1, 1, 1, 1, 1, -4, 0, 0) \\ I(1, 1, 1, 1, 1, 1, 1, 1, 0, 0, -3), \quad I(1, 1, 1, 1, 1, 1, 1, 1, 0, -1, -2), \quad I(1, 1, 1, 1, 1, 1, 1, 1, 0, -2, -1) \\ I(1, 1, 1, 1, 1, 1, 1, 1, 0, -3, 0), \quad I(1, 1, 1, 1, 1, 1, 1, 1, 0, -1, -2), \quad I(1, 1, 1, 1, 1, 1, 1, 1, -1, -1, -1) \\ I(1, 1, 1, 1, 1, 1, 1, 1, -1, -2, 0), \quad I(1, 1, 1, 1, 1, 1, 1, 1, -2, 0, -1), \quad I(1, 1, 1, 1, 1, 1, 1, 1, -2, -1, 0) \\ I(1, 1, 1, 1, 1, 1, 1, 1, -3, 0, 0), \quad I(1, 1, 1, 1, 1, 1, 1, 1, 0, 0, -2), \quad I(1, 1, 1, 1, 1, 1, 1, 1, 0, -1, -1) \\ I(1, 1, 1, 1, 1, 1, 1, 1, 1, 0, -2, 0), \quad I(1, 1, 1, 1, 1, 1, 1, 1, 1, -1, 0, -1), \quad I(1, 1, 1, 1, 1, 1, 1, 1, -1, -1, 0) \\ I(1, 1, 1, 1, 1, 1, 1, 1, 1, 0, -1), \quad I(1, 1, 1, 1, 1, 1, 1, 1, 1, 0, -1, 0) \}$$

can be downloaded from (268 MB compressed / 790 MB uncompressed)

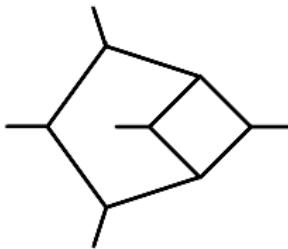
https://github.com/yzhphy/hexagonbox_reduction/releases/download/1.0.0/hexagon_box_degree_4_Final.zip

- Our results agree with fully numerical results from FIRE5 C++
(6 hours per point).

[A. Smirnov, CPC 189(2015)182]

Conclusions

- New formalism for IBP reductions. Main ideas: **cuts**, IBP identities from syzygies, total pivoting, rational reconstruction
- Obtained the **fully analytic** IBP reductions of



with numerator insertions up to degree 4 in the z_i .

- Powerful framework. IBP reductions for further $2 \rightarrow 3$ two-loop processes **seem well within reach**.